CS 5150/6150: Assignment 2 Due: Sep 12, 2011

This assignment has 6 questions, for a total of 100 points and 0 bonus points. Unless otherwise specified, complete and reasoned arguments will be expected for all answers.

Solving a recurrence means that you provide a bound of the form T(n) = O(f(n)) for a specific f. Tight bounds get full credit: for example, if the recurrence is T(n) = 2T(n/2) + cn, the answer $T(n) = O(n^2)$ is correct but not precise enough, and you will not get full credit for it.

- (a) [4] T(n) = 3T(n/2) + cn
- (b) [4] $T(n) = 4T(n/4) + n\log n$
- (c) [4] $T(n) = 5T(n/8) + \sqrt{n}$
- (d) [4] $T(n) = 2T(n/2) + \frac{n}{\log n}$
- (e) [4] $T(n) = \frac{T(n-1)}{T(n-2)}, T(0) = 1, T(1) = 2.$
- Question 2: Finding a local minimum......[10]
 Please solve problem 11 from http://compgeom.cs.uiuc.edu/~jeffe/teaching/algorithms/notes/
 01-recursion.pdf

A line is *visible* if there is some x_0 for which it is uppermost. Output the set of visible lines as efficiently as possible.

- - (a) Implement an algorithm for *k*-selection.

- (b) Experiment with different choices for finding a pivot. Specifically, use different groupings of the input and report on how this affects the performance
- (c) How does this algorithm compare to sorting and then taking the k^{th} element in practice ? Provide concrete empirical evidence in support of your answer.
- (d) How do random pivots compare to computing deterministic pivots ? Again, experiment with different random pivots and provide concrete empirical evidence in support of your answer.