

# ALGORITHM ANALYSIS

cs2420 | Introduction to Algorithms and Data Structures | Spring 2015

**administrivia...**

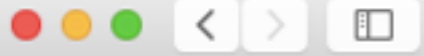
-assignment 2 is due Friday at midnight

-note change in due date, and time

-tutoring experiment

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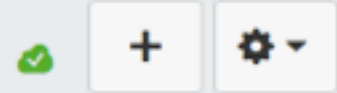
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# ARE YOU HERE?

- 1) yes
- 2) no

**last time...**



**generics**

-we always see `<>` associated with `ArrayList`...

```
ArrayList<Shape> list = new ArrayList<Shape>();
```

-`ArrayList` is a **generic class** — we can create any version of it that we want

-**generic programming:** algorithms are written in terms of types to-be-specified-later

-algorithms instantiated when needed for specific types defined by parameters

-here's what the code actually looks like:

```
public class ArrayList<T> {  
    T storage[];  
    int capacity, numItems;  
  
    public void add(T item)  
    { ... }  
}
```

-the placeholder `T` is replaced with the real type when you instantiate an `ArrayList` with `<>`

-`T` can be used as a type anywhere in `ArrayList` class

-generics allow for type-checking at compile time instead of run-time

-can detect type mismatch **BEFORE** your code runs

## BEFORE GENERICS:

```
ArrayList l;  
l.add(new String("hi"));  
Shape i = (Shape)l.get(0); // crash
```

## ALTERNATIVE:

```
ArrayList<String> l;  
l.add(new String("hi"));  
Shape i = (Shape)l.get(0); // compile error
```

COMPILE-TIME ERRORS ARE ALWAYS BETTER THAN RUN-TIME!

-static methods can have their own generic types

-declare the generic type before the return type:

```
public static <T> boolean doWork(...) {...}
```

-we can refer to T as a type within that method only!

-example:

```
public static <T> boolean contains(T[] array, T item)
{
    for(int i=0; i < array.length; i++)
        if(array[i].equals(item))
            return true;

    return false;
}
```

**today...**

-algorithm analysis

-complexity growth rate

-big-O notation



# algorithm analysis

- correctness is only half the battle
- programs are expected to terminate in a reasonable amount of time
- running time of a program is strongly correlated to the choice of algorithms used in problem solving
- how much time and space does an algorithm require?**

**example...**

# finding a word in a dictionary

## ALGORITHM 1:

- 1) start on the **first** page, **first** entry
- 2) if word not found, move to the next entry
- 3) if very end of dictionary reached, word not found

## IS THIS ALGORITHM CORRECT?

- 1) yes
- 2) no

## CAN WE DO BETTER?

- 1) yes
- 2) no

# finding a word in a dictionary

## ALGORITHM 2:

1) guess which page the entry is on

2) did we go too far?

- go back some pages

3) did we not go far enough?

- go forward some pages

4) continue narrowing

WHAT DOES THIS ALGORITHM ASSUME ABOUT THE DICTIONARY?

## -ALGORITHM 1: LINEAR SEARCH

-running time directly related to size of dictionary

*-assume 180K words, and 0.25s to check one word*

*-12 hours to complete!*

## -ALGORITHM 2: BINARY SEARCH

-more like what humans do

*-4 seconds to complete!*

**TIME \* 2 → O(N)**

**TIME + 0.25 → O(logN)**

-what if the dictionary doubles?

### ALGORITHM 1 RUN-TIME?

- 1) time \* 0.5
- 2) time \* 2
- 3) time + 0.25
- 4) time + 10

### ALGORITHM 2 RUN-TIME?

- 1) time \* 0.5
- 2) time \* 2
- 3) time + 0.25
- 4) time + 10

# a note on logarithms

-a **logarithm** is an exponent indicating the power to which a base is raised to produce a given number

$$\log_B N = X$$

$$B^X = N$$

HOW MANY BITS DOES IT TAKE TO REPRESENT A NUMBER?

-by default the base is 2 ...we'll come back to this

-the logarithm grows slowly

$$9 < \log 1000 < 10$$

$$19 < \log 1,000,000 < 20$$

$$29 < \log 1,000,000,000 < 30$$

- **$N \log N$**  is closer to  **$N$**  than  **$N^2$**

why is binary search  $O(\log N)$ ?

why is the default base 2?



# finding a word in a dictionary

-binary search will **always** win for large dictionaries  
-as **N** increases, the gap between the algorithms becomes larger

-linear search has linear growth rate

-graph is a \_\_\_\_\_ line

-run time for  **$N = T$**  units of time

-run time for  **$2N = 2T$**  units of time

- 1) exponential
- 2) straight
- 3) negative-slope

-binary search has logarithmic growth rate

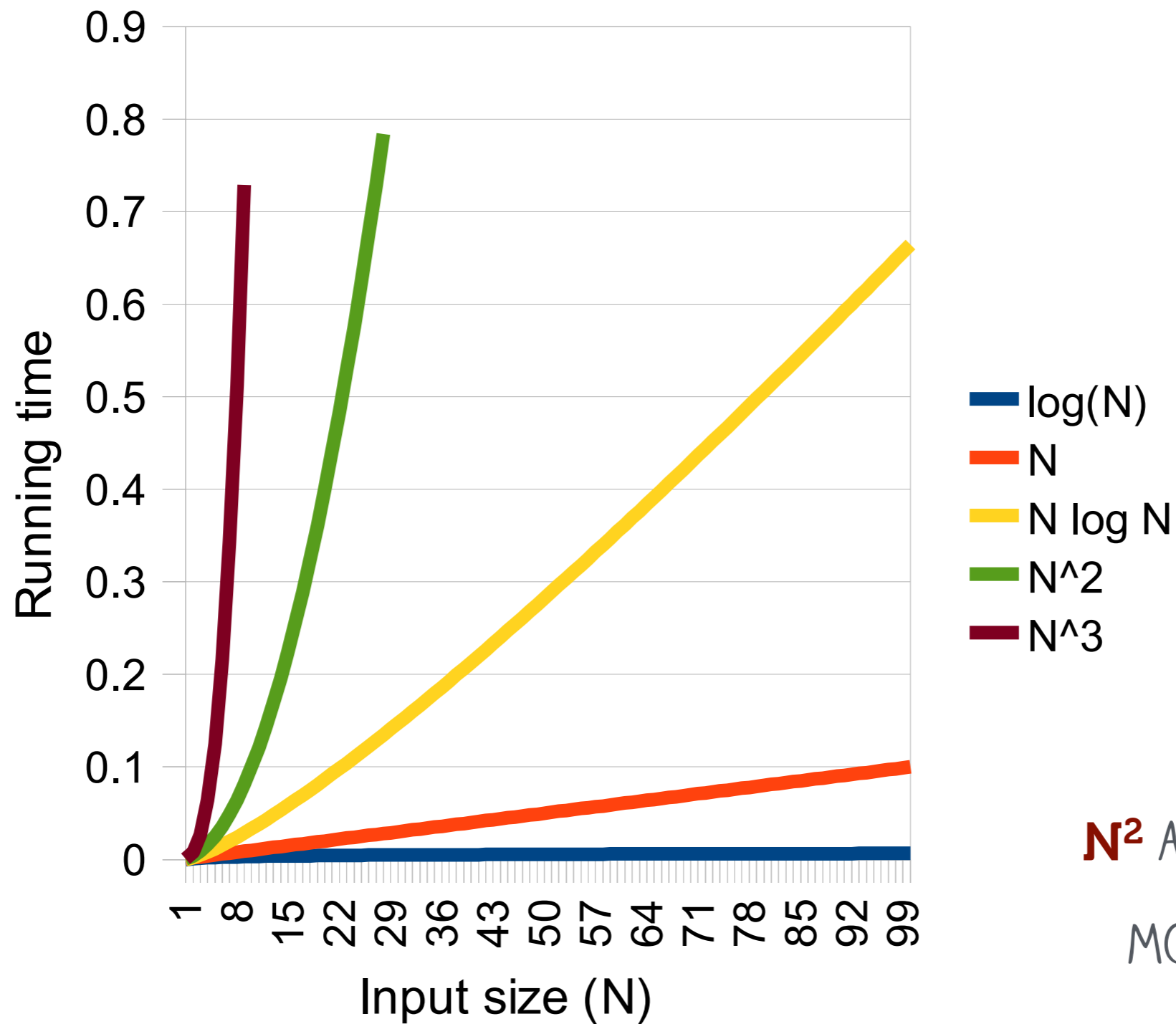
-run time for  **$N = T$**  units of time

-run time for  **$2N = T + \_$**  units of time

- 1) 1
- 2) 2
- 3) 10

**growth rate**

# typical run-time complexities



**N<sup>2</sup>** AND **N<sup>3</sup>** ARE TYPICALLY NOT ACCEPTABLE FOR MODERATE INPUT SIZES!

-knowing that  $F(N) < G(N)$  for a particular  $N$  is not very useful

-instead, we measure the functions' growth rates

-for sufficiently large  $N$ , a function's growth rate is determined by its dominate term

$10N^2 + 40N + 760$   $\longrightarrow$  WHAT IS THE DOMINATE TERM?

<b>c</b>	constant
<b>log N</b>	logarithmic
<b>N</b>	linear
<b>N log N</b>	linearithmic
<b>N<sup>2</sup></b>	quadratic
<b>N<sup>3</sup></b>	cubic

$\downarrow$  increasing growth rate

# how to get log growth?

-how many bits are needed to represent **N** consecutive integers?

-starting at **x=1**, how many iterations of **x\*2** before **x>=N**?

-the *repeated doubling* principle

-starting at **x=N**, how many iterations of **x/2** before **x<=1**?

-the *repeated halving* principle

# big-O notation

-**big-O notation** (**O**) is used to capture the dominate term in an algorithm

-assuming large **N**!

-for example, the running time of a quadratic algorithm is **N<sup>2</sup>** is specified **O(N<sup>2</sup>)**

-pronounced “order N squared”

-this notation allows us to establish a relative order among algorithms

-**O(N log N)** is better than **O(N<sup>2</sup>)**

# what's code got to do, got to do with it...

- $O(N^2)$  and  $O(N^3)$  are impractical for most  $N$
- clever programming tricks **CANNOT** make an inefficient algorithm fast
  - a poorly coded linear algorithm trumps a quadratic algorithm in a highly efficient machine language

## TAKE AWAY:

OPTIMIZING THE ALGORITHM (OR CHOOSING THE BEST ONE) WILL GET YOU MUCH FURTHER THAN OPTIMIZING THE CODE



# worst, average, best

- worst-case** is a guarantee on all inputs — it will never be worse than this
- average-case** is the common case, measured over all possible inputs
  - this is the most useful!
- best-case** is the absolute fastest that an algorithm can terminate
  - we don't care about this because it rarely happens

**example...**

# finding the maximum item in an array

## ALGORITHM?

- 1) initialize `max` to the first element
- 2) scan through each item in the array
  - if the item is greater than `max`, update `max`

## WHAT IS THE BIG-O COMPLEXITY OF THIS ALGORITHM?

- 1) **c**
- 2)  **$\log N$**
- 3)  **$N$**
- 4)  **$N \log N$**
- 5)  **$N^2$**
- 6)  **$N^3$**

# finding the smallest difference

## ALGORITHM?

```
diff = MAX_INTEGER;
for(int i=0; i<array.length-1; i++)
{
    num1 = array[i];
    for(int j=i+1; j<array.length; j++)
    {
        num2 = array[j];
        if (abs(num1-num2) < diff)
            diff = abs(num1-num2);
    }
}
return diff;
```

- 1) **c**
- 2) **log N**
- 3) **N**
- 4) **N log N**
- 5) **N<sup>2</sup>**
- 6) **N<sup>3</sup>**

WHAT IS THE BIG-O COMPLEXITY OF THIS ALGORITHM?

**next time...**

-reading

-chapters 5 & 6

-homework

-assignment 2 due Friday at 11:59pm

*-must complete with a partner!*