

INTRO TO SORTING

cs2420 | Introduction to Algorithms and Data Structures | Spring 2015

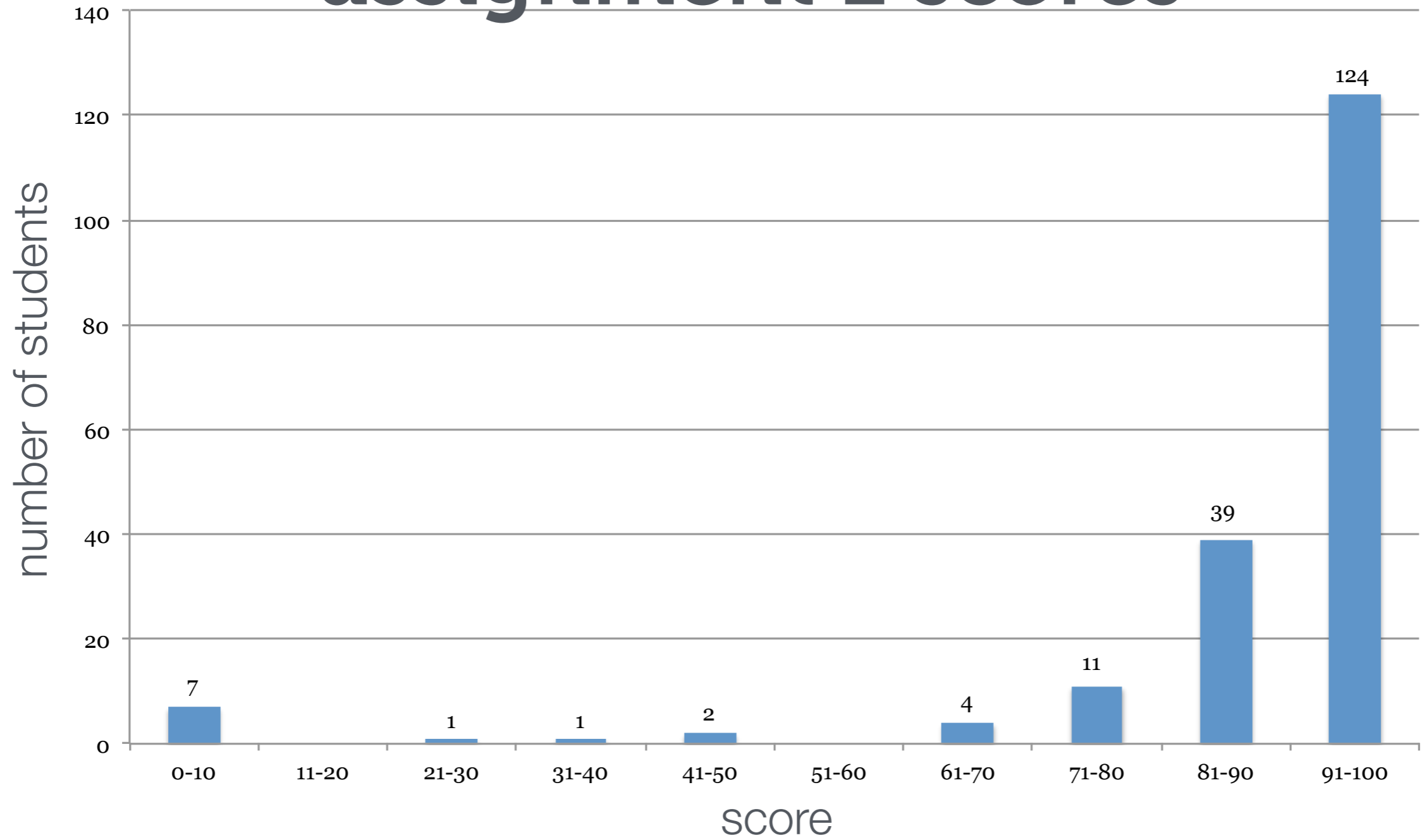
administrivia...

-assignment 3 is due Thursday at midnight (11:59pm)

-tutoring

-Doodle polls up on the website

assignment 1 scores



CLICKERS!!!
...it's gonna work!



ARE YOU HERE?

A) yes

B) no

last time...

-a Collection is a data structure that holds items

-very unspecific as to how the items are held

-ie. the data structure

-supports various operations:

-add, remove, contains, ...

-examples:

-ArrayList

-PriorityQueue

-LinkedList

-TreeSet

WHAT IF WE USE AN ARRAY UNDER THE HOOD?

add

```
int[] data = new int[6];  
data.add(5);  
data.add(17);  
data.add(9);  
data.add(12);  
data.add(1);  
data.add(33);
```

5	17	9	12	1	33
---	----	---	----	---	----

```
data.add(22);
```

NOW WHAT???

WHAT IS THE COMPLEXITY OF add?

- A) **c**
- B) **log N**
- C) **N**
- D) **N log N**
- E) **N²**
- F) **N³**

grow

data →

5	17	9	12	1	33
---	----	---	----	---	----

```
tmp = new int[data.length*2];
```

tmp →

--	--	--	--	--	--	--	--	--	--	--	--

copy all from data to tmp

tmp →

5	17	9	12	1	33						
---	----	---	----	---	----	--	--	--	--	--	--

```
data = tmp;
```

data →

5	17	9	12	1	33						
---	----	---	----	---	----	--	--	--	--	--	--

- A) **c**
- B) **log N**
- C) **N**
- D) **N log N**
- E) **N²**
- F) **N³**


WHAT IS THE COMPLEXITY OF GROWING?

remove

5	17	9	12	1	33
---	----	---	----	---	----

`data.remove(9);`

5	17		12	1	33
---	----	--	----	---	----



The diagram shows three curved arrows pointing from the right towards the left, indicating that elements from index 3 to 5 are shifted one position to the right to fill the gap at index 2.

5	17	12	1	33	
---	----	----	---	----	--

WHAT IS THE COMPLEXITY OF `remove`?

- A) **c**
- B) **log N**
- C) **N**
- D) **N log N**
- E) **N²**
- F) **N³**

Iterator

-an `Iterator` is specific to a data structure, and knows how to traverse the structure

- `hasNext`: determines if iteration is complete

- `next`: gets the next item

- `remove`: removes the last seen item

-internally, keeps track of where the next item is (as well as other state)

-actually points to *between* items

today...

-why sort?

-selection sort

-insertion sort

why sort?

- sorting is a fundamental application in computing
 - one of the most intensively studied and important operations
- most data is useless unless it is in some kind of order
- for any given problem, or specific goal isn't necessarily sorting... but we often need to sort to efficiently solve problems
 - computer graphics
 - look-up tables
 - games

- sorting algorithms that are easy to understand (and implement) run in **quadratic time**
- more complicated algorithms cut it to **$O(N \log N)$**
 - implementation details are critical to attaining this bound!
- for very specific types of data we can actually do better
 - but we won't study these algorithms extensively

WITHOUT THINKING TOO HARD, HOW
CAN WE SORT ANY ARRAY OF ITEMS?

selection sort

the simplest sorting algorithm

selection sort

- 1) find the minimum item in the unsorted part of the array
- 2) swap it with the first item in the unsorted part of the array
- 3) repeat steps 1 and 2 to sort the remainder of the array

WHAT DOES THIS LOOK LIKE?

```

void selectionSort(int[] arr)
{
    for(int i=0; i < arr.length-1; i++)
    {
        min = i; —— LAST ITEM IN SORTED PART OF ARRAY
        for(int j=i+1; j < arr.length; j++)
            if (arr[j] < arr[min])
                min = j;

        temp = arr[i];
        arr[i] = arr[min];
        arr[min] = temp; SWAP ITEMS
    }
}

```

**LOOK FOR ITEM LESS THAN
THOSE IN SORTED PART OF ARRAY**

WHAT IS THE COMPLEXITY OF SELECTION SORT?

L1 `for(int i=0; i < arr.length-1; i++)`

L2 `for(int j=i+1; j < arr.length; j++)`


```
void selectionSort(int[] arr)
{
    for(int i=0; i < arr.length-1; i++)
    {
        min = i;
        for(int j=i+1; j < arr.length; j++)
            if (arr[j] < arr[min])
                min = j;

        temp = arr[i];
        arr[i] = arr[min];
        arr[min] = temp;
    }
}
```

- A) **c**
- B) **log N**
- C) **N**
- D) **N log N**
- E) **N²**
- F) **N³**

WHAT IS THE BEST-CASE COMPLEXITY OF SELECTION SORT?

insertion sort

good for small **N**

insertion sort

- 1) the first array item is the sorted portion of the array
- 2) take the second item and insert it in the sorted portion
- 3) repeat steps 1 and 2 to sort the remainder of the array

WHAT DOES THIS LOOK LIKE?

```

void insertionSort(int[] arr)
{
    for(int i=1; i < arr.length; i++)
    {
        index = arr[i]; —— ITEM TO BE INSERTED
        j = i;
        while(j>0 && arr[j-1]>index)
        {
            arr[j] = arr[j-1];
            j--;
        }
        arr[j] = index; —— INSERT ITEM
    }
}

```

UNTIL THE INSERTION
POSITION IS FOUND,
SHIFT SORTED ITEMS



WHAT IS THE COMPLEXITY OF INSERTION SORT?

unsortedness

-requires a measure of *unsortedness* for array

-**inversion**: a pair of array items that are out of order

45	-3	9	76	11	-8	0
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HOW MANY INVERSIONS ARE THERE?

-sorting efficiency depends on how many inversions are removed per step

insertion sort complexity

each swap to the left removes one inversion...

...we must visit each item at least once (**N**)...

...and we must undo **I** inversions

45	-3	9	76	11	-8	0
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SWAP REMOVES ONE INVERSION

insertion sort is **$O(N+I)$**

HOW DO WE FIGURE OUT WHAT **I** IS?

next time...

-reading

-chapters 8.1 - 8.4

-homework

-assignment 3 due on Thursday